

Similarity Search

CSE545 - Spring 2020
Stony Brook University

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$A \cap B$

Big Data Analytics, The Class

Goal: Generalizations
A model or summarization of the data.

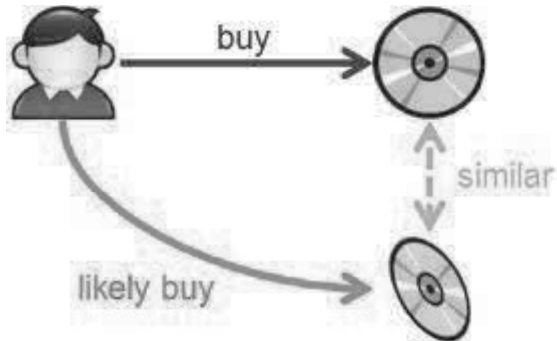
Data Frameworks

Algorithms and Analyses

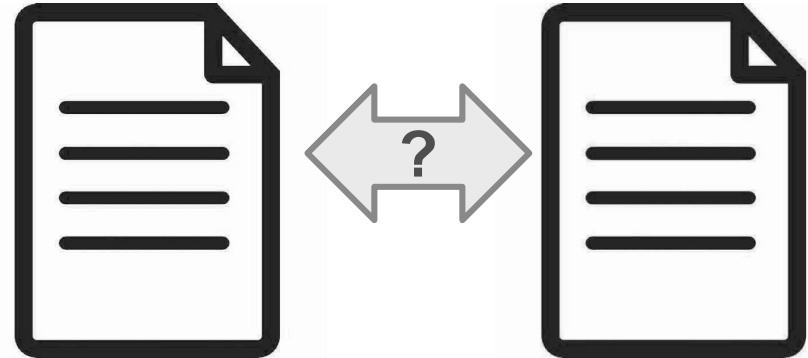
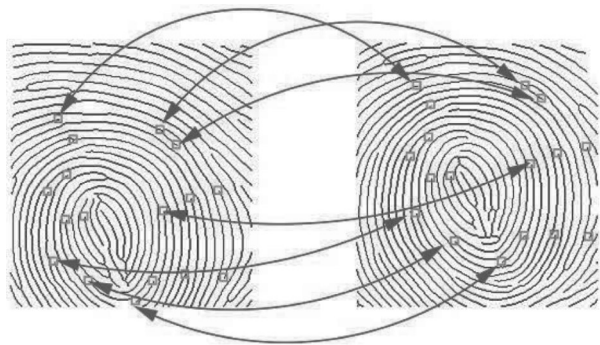
*Hadoop File System ✓ Spark ✓
Streaming ✓
MapReduce ✓ Tensorflow ✓*

*Similarity Search
Hypothesis Testing
Link Analysis
Recommendation Systems
Deep Learning*

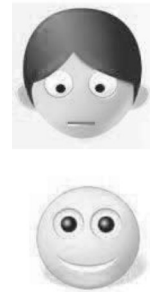
Finding Similar Items



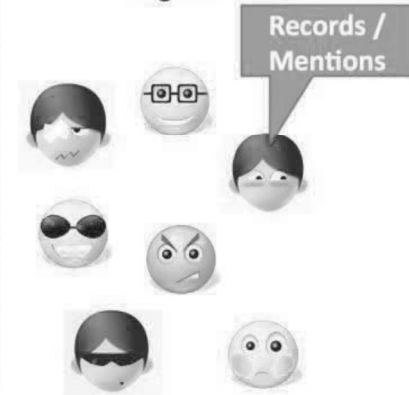
(<http://blog.soton.ac.uk/hive/2012/05/10/recommendation-system-of-hive/>)



Real World



Digital World



(<http://www.datacommunitydc.org/blog/2013/08/entity-resolution-for-big-data>)

Finding Similar Items: Topics

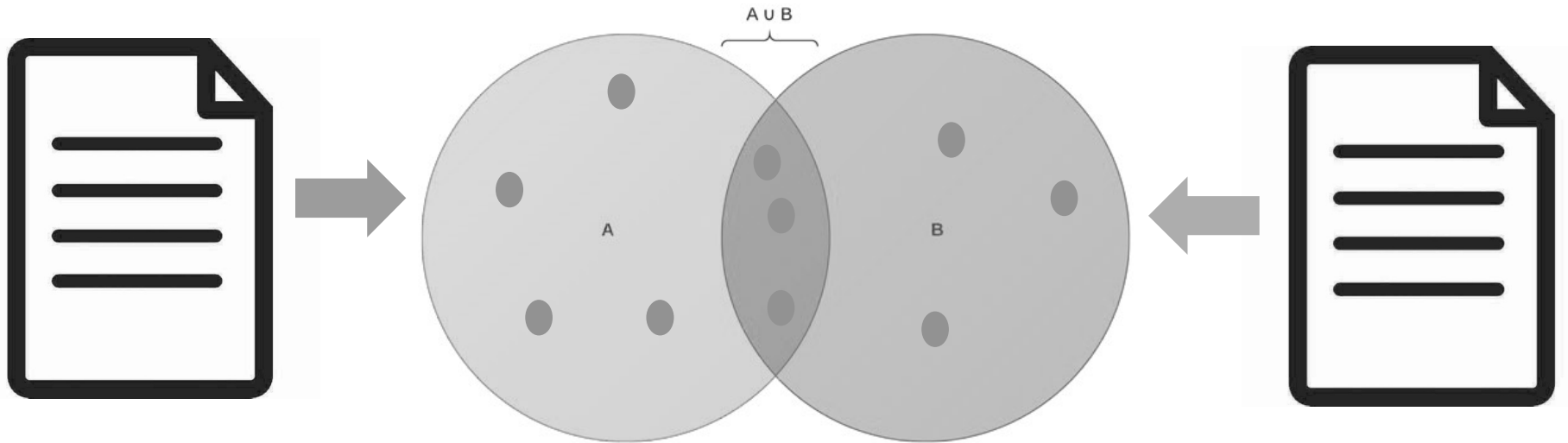
- Shingling
- Minhashing
- Locality-sensitive hashing
- Distance Metrics

Document Similarity

Challenge: How to represent the document in a way that can be efficiently encoded and compared?

Shingles

Goal: Convert documents to sets



Shingles

Goal: Convert documents to sets



k -shingles (aka “character n -grams”)
- sequence of k characters



E.g. $k=2$ doc=“abcdabd”
 $\text{singles}(\text{doc}, 2) = \{\text{ab}, \text{bc}, \text{cd}, \text{da}, \text{bd}\}$

Shingles

Goal: Convert documents to sets



k -shingles (aka “character n -grams”)
- sequence of k characters



E.g. $k=2$ doc=“abcdabd”
 $\text{singles}(\text{doc}, 2) = \{\text{ab}, \text{bc}, \text{cd}, \text{da}, \text{bd}\}$

- Similar documents have many common shingles
- Changing words or order has minimal effect.
- In practice use $5 < k < 10$

Shingles

Goal: Convert documents to sets



Large enough that any given shingle appearing a document is highly unlikely (e.g. $< .1\%$ chance)

Can hash large shingles to smaller (e.g. 9-shingles into 4 bytes)

Can also use words (aka n-grams).

- Similar documents have many common shingles
- Changing words or order has minimal effect.
- **In practice use $5 < k < 10$**

Shingles

Problem: Even if hashing, sets of shingles are large
(e.g. 4 bytes \Rightarrow 4x the size of the document).

Minhashing

Goal: Convert sets to shorter ids, signatures

Minhashing

Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix, X :

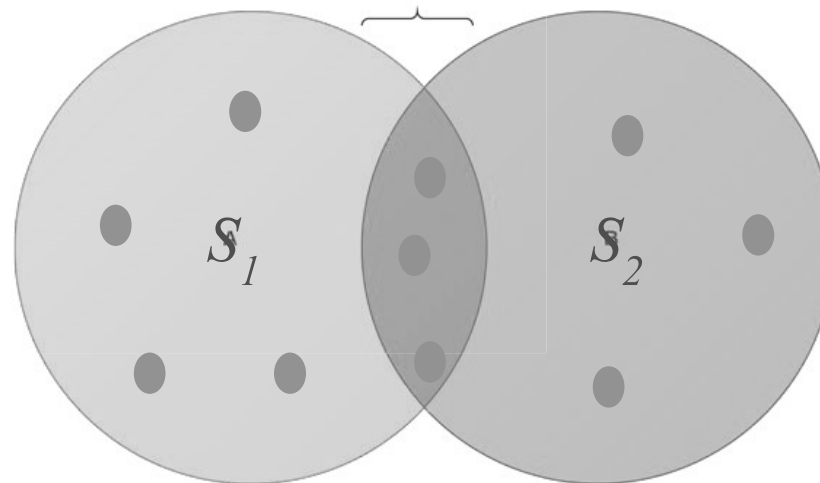
<i>Element</i>	S_1	S_2	S_3	S_4
<i>a</i>	1	0	0	1	
<i>b</i>	0	0	1	0	
<i>c</i>	0	1	0	1	
<i>d</i>	1	0	1	1	
<i>e</i>	0	0	1	0	

(Leskovec et al., 2014; <http://www.mmids.org/>)

often very sparse! (lots of zeros)

Jaccard Similarity:

$$\text{sim}(S_1, S_2) = \frac{S_1 \cap S_2}{S_1 \cup S_2}$$



Minhashing

Characteristic Matrix:

	S_1	S_2
ab	1	1
bc	0	1
de	1	0
ah	1	1
ha	0	0
ed	1	1
ca	0	1

Jaccard Similarity:

$$\text{sim}(S_1, S_2) = \frac{S_1 \cap S_2}{S_1 \cup S_2}$$

Minhashing

Characteristic Matrix:

	S_1	S_2	
ab	1	1	**
bc	0	1	*
de	1	0	*
ah	1	1	**
ha	0	0	
ed	1	1	**
ca	0	1	*

Jaccard Similarity:

$$sim(S_1, S_2) = \frac{S_1 \cap S_2}{S_1 \cup S_2}$$

Minhashing

Characteristic Matrix:

	S_1	S_2	
ab	1	1	**
bc	0	1	*
de	1	0	*
ah	1	1	**
ha	0	0	
ed	1	1	**
ca	0	1	*

Jaccard Similarity:

$$sim(S_1, S_2) = \frac{S_1 \cap S_2}{S_1 \cup S_2}$$

$$sim(S_1, S_2) = 3 / 6$$

both have / # at least one has

Minhashing

Problem: Even if hashing shingle contents,
sets of shingles are large

e.g. 4 byte integer per shingle: assume all unique shingles,
=> 4x the size of the document

(since there are as many shingles as characters and 1byte per char).

Minhashing

Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix: X

	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

(Leskovec et al., 2014; <http://www.mmds.org/>)

Minhashing

Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix: X

	S_1	S_2	S_3	S_4
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de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0


Approximate Approach:

- 1) Instead of keeping whole characteristic matrix, just keep first row where 1 is encountered.
- 2) Shuffle and repeat to get a “signature” for each set.

Minhashing

Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix: X



	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
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ed	1	0	1	0
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ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

(Leskovec et al., 2014; <http://www.mmhds.org/>)

Approximate Approach:

1) Instead of keeping whole characteristic matrix, just keep first row where 1 is encountered.

2) Shuffle and repeat to get a “signature”.

	S_1	S_2	S_3	S_4
ah	0	1	0	1
ca	1	0	1	0
ed	1	0	1	0
de	0	1	0	1
ab	1	0	1	0
bc	1	0	0	1

...

Minhashing

Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix: X

	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

(Leskovec at al., 2014; <http://www.mmids.org/>)

Approximate Approach:

1) Instead of keeping whole characteristic matrix, just keep first row where 1 is encountered.

2) Shuffle and repeat to get a “signature”.

	S_1	S_2	S_3	S_4
ah	0	1	0	1
ca	1	0	1	0
ed	1	0	1	0
de	0	1	0	1
ab	1	0	1	0
bc	1	0	0	1

signatures

S_1	S_2	S_3	S_4
1	3	1	2
2	1	2	1
...

■ ■ ■

Minhashing

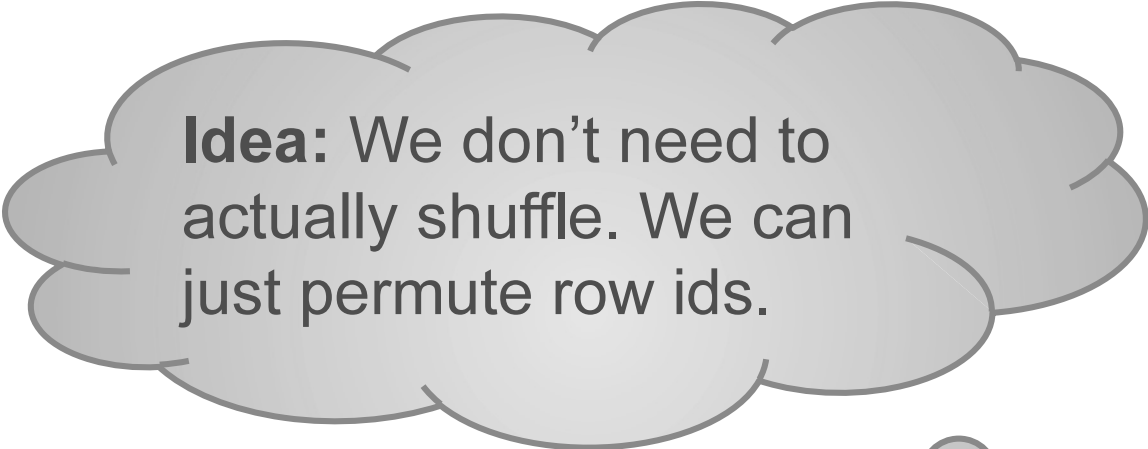
Goal: Convert sets to shorter ids, “signatures”

Characteristic Matrix: X

	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

Approximate Approach:

- 1) Instead of keeping whole characteristic matrix, just keep first row where 1 is encountered.
- 2) Shuffle and repeat to get a “signature” for each set.



Idea: We don't need to actually shuffle. We can just permute row ids.

Minhashing

Characteristic Matrix:

	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

Minhashing

Characteristic Matrix:

	S_1	S_2	S_3	S_4
ab	1	0	1	0
bc	1	0	0	1
de	0	1	0	1
ah	0	1	0	1
ha	0	1	0	1
ed	1	0	1	0
ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

permuted order
1 ha
2 ed
3 ab
4 bc
5 ca
6 ah
7 de

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

permuted order
1 ha
2 ed
3 ab
4 bc
5 ca
6 ah
7 de

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

permuted order
1 ha
2 ed
3 ab
4 bc
5 ca
6 ah
7 de

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

$$h(S_1) = \text{ed} \quad \# \text{permuted row 2}$$

$$h(S_2) = \text{ha} \quad \# \text{permuted row 1}$$

$$h(S_3) =$$

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

permuted order
1 ha
2 ed
3 ab
4 bc
5 ca
6 ah
7 de

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

$$h(S_1) = ed \text{ \#permuted row 2}$$

$$h(S_2) = ha \text{ \#permuted row 1}$$

$$h(S_3) = ed \text{ \#permuted row 2}$$

$$h(S_4) =$$

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

permuted order
1 ha
2 ed
3 ab
4 bc
5 ca
6 ah
7 de

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to first row where set appears.

$$h(S_1) = \text{ed} \quad \# \text{permuted row 2}$$

$$h(S_2) = \text{ha} \quad \# \text{permuted row 1}$$

$$h(S_3) = \text{ed} \quad \# \text{permuted row 2}$$

$$h(S_4) = \text{ha} \quad \# \text{permuted row 1}$$

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

(Leskovec et al., 2014; <http://www.mmhds.org/>)

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1

$h_1(S_1) = ed$ #permuted row 2

$h_1(S_2) = ha$ #permuted row 1

$h_1(S_3) = ed$ #permuted row 2

$h_1(S_4) = ha$ #permuted row 1

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

(Leskovec et al., 2014; <http://www.mmids.org/>)

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1

$$h_1(S_1) = \text{ed} \quad \# \text{permuted row}$$

2

$$h_1(S_2) = \text{ha} \quad \# \text{permuted row}$$

1

$$h(S_3) = \text{ed} \quad \# \text{permuted row}$$

Minhashing

Characteristic Matrix:

		S_1	S_2	S_3	S_4
3	ab	1	0	1	0
4	bc	1	0	0	1
7	de	0	1	0	1
6	ah	0	1	0	1
1	ha	0	1	0	1
2	ed	1	0	1	0
5	ca	1	0	1	0

(Leskovec at al., 2014; <http://www.mmids.org/>)

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- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1

$$h_1(S_1) = ed \text{ \#permutated row}$$

2

$$h_1(S_2) = ha \text{ \#permutated row}$$

1

$$h(S_3) = ed \text{ \#permutated row}$$

Minhashing

Characteristic Matrix:

			S_1	S_2	S_3	S_4
4	3	ab	1	0	1	0
2	4	bc	1	0	0	1
1	7	de	0	1	0	1
3	6	ah	0	1	0	1
6	1	ha	0	1	0	1
7	2	ed	1	0	1	0
5	5	ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2				

Minhashing

Characteristic Matrix:

			S_1	S_2	S_3	S_4
4	3	ab	1	0	1	0
2	4	bc	1	0	0	1
1	7	de	0	1	0	1
3	6	ah	0	1	0	1
6	1	ha	0	1	0	1
7	2	ed	1	0	1	0
5	5	ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	1	0
3	2	4	bc	1	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3				

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	1	0
3	2	4	bc	1	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	1	0
3	2	4	bc	1	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

(Leskovec et al., 2014; <http://www.mmms.org/>)

Minhash function: h

- Based on permutation of rows in the characteristic matrix, h maps sets to rows.

Signature matrix: M

- Record first row where each set had a 1 in the given permutation

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2
...				
...				

Minhashing

Characteristic Matrix:

Property of signature matrix:
 The probability for any h_i (i.e. any row), that $h_i(S_1) = h_i(S_2)$ is the same as $\text{Sim}(S_1, S_2)$

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	1	0
3	2	4	bc	1	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2
...				
...				

(Leskovec et al., 2014; <http://www.mmds.org/>)

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	1	0
3	2	4	bc	1	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

Property of signature matrix:
 The probability for any h_i (i.e. any row), that $h_i(S_1) = h_i(S_2)$ is the same as $\text{Sim}(S_1, S_2)$

Thus, similarity of signatures S_1, S_2 is the fraction of minhash functions (i.e. rows) in which they agree.

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2
...				
...				

(Leskovec at al., 2014; <http://www.mmds.org/>)

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	0	0
3	2	4	an	0	1	0	0
6	3	6	an	0	1	0	0
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

Property of signature matrix:

The probability for any h_i (i.e. any row), that $h_i(S_1) = h_i(S_2)$ is the same as $\text{Sim}(S_1, S_2)$

Thus, similarity of signatures S_1, S_2 is the fraction of minhash functions (i.e. rows) in which they agree.

Estimate with a random sample of permutations (i.e. ~100)

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2
...				
...				

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	1	0	0	0
3	2	4	ba	0	1	0	0
6	3	6	an	0	0	1	0
2	6	1	ha	0	1	0	1
5	7	2	ed	1	0	1	0
4	5	5	ca	1	0	1	0

Property of signature matrix:

The probability for any h_i (i.e. any row), that $h_i(S_1) = h_i(S_2)$ is the same as $\text{Sim}(S_1, S_2)$

Thus, similarity of signatures S_1, S_2 is the fraction of minhash functions (i.e. rows) in which they agree.

Estimate with a random sample of permutations (i.e. ~100)

	S_1	S_2	S_3	S_4
h_1	2	1	2	1
h_2	2	1	4	1
h_3	1	2	1	2

Estimated $\text{Sim}(S_1, S_3) = \text{agree} / \text{all} = 2/3$

Minhashing

Characteristic Matrix:

				S_1	S_2	S_3	S_4
1	4	3	ab	<u>1</u>	0	<u>1</u>	0
3	2	4	bc	<u>1</u>	0	0	1
7	1	7	de	0	1	0	1
6	3	6	ah	0	1	0	1
2	6	1	ha	0	1	0	1
5	7	2	ed	<u>1</u>	0	<u>1</u>	0
4	5	5	ca	<u>1</u>	0	<u>1</u>	0

Property of signature matrix:
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Minhashing

Characteristic Matrix:

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- Can't reasonably do permutations (huge space)
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Solution: Use “random” hash functions.

- Setup:
 - Pick ~ 100 hash functions, hashes
 - Store $M[i][s] = \text{a potential minimum } h_i(r)$
#initialized to infinity (num hashes x num sets)

Minhashing

Solution: Use “random” hash functions.

Setup:

```
hashes = [getHfunc(i) for i in rand(1, num=100)]
```

#100 hash functions, seeded random

```
for i in hashes: for s in sets:
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```
    Sig[i][s] = np.inf #represents a potential minimum  $h_i(r)$  ; initially infinity
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Algorithm (“efficient minhashing”):

```
for r in rows of cm: #cm is characteristic matrix  
    compute  $h_i(r)$  for all i in hashes #precompute 100 values  
    for each set s in sets: #columns of cm  
        if cm[r][s] == 1:  
            for i in hashes: #check which hash produces smallest value  
                if  $h_i(r) < \text{Sig}[i][s]$ : Sig[i][s] =  $h_i(r)$ 
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Algorithm (“efficient minhashing”) without charact matrix:


```
for feat in shins: #shins is all unique shingles  
    compute  $h_i(\text{feat})$  for all i in hashes #precompute 100 values  
    for each set s in sets: #sets is list of shingle sets  
        if feat in s:  
            for i in hashes: #check which hash produces smallest value  
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Minhashing

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New Problem: Even if the size of signatures are small, it can be computationally expensive to find similar pairs.

E.g. 1m documents; 1,000,000 choose 2 = 500,000,000,000 pairs!

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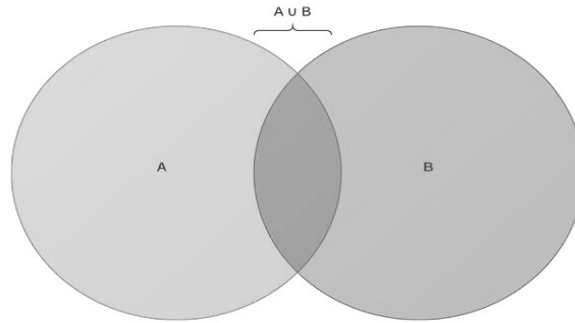
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(1m documents isn't even "big data")

Document Similarity



Duplicate web pages (useful for ranking

Plagiarism

Cluster News Articles

Anything similar to documents: movie/music/art tastes, product characteristics

COVID-19 Report matching

Locality-Sensitive Hashing

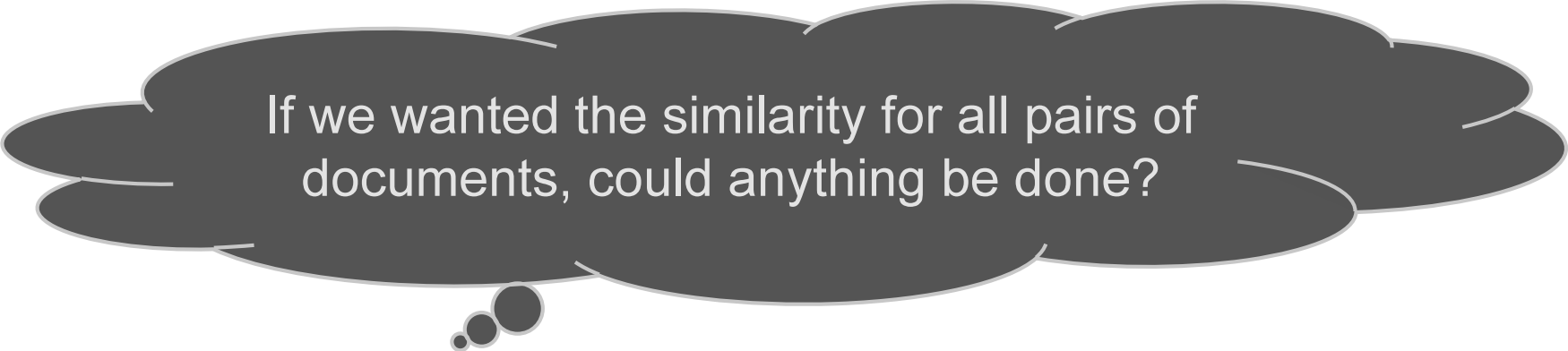
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Candidate pairs: pairs of elements to be evaluated for similarity.

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If we wanted the similarity for all pairs of documents, could anything be done?

Locality-Sensitive Hashing

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Approach: Hash multiple times over subsets of data: similar items are likely in the same bucket once.

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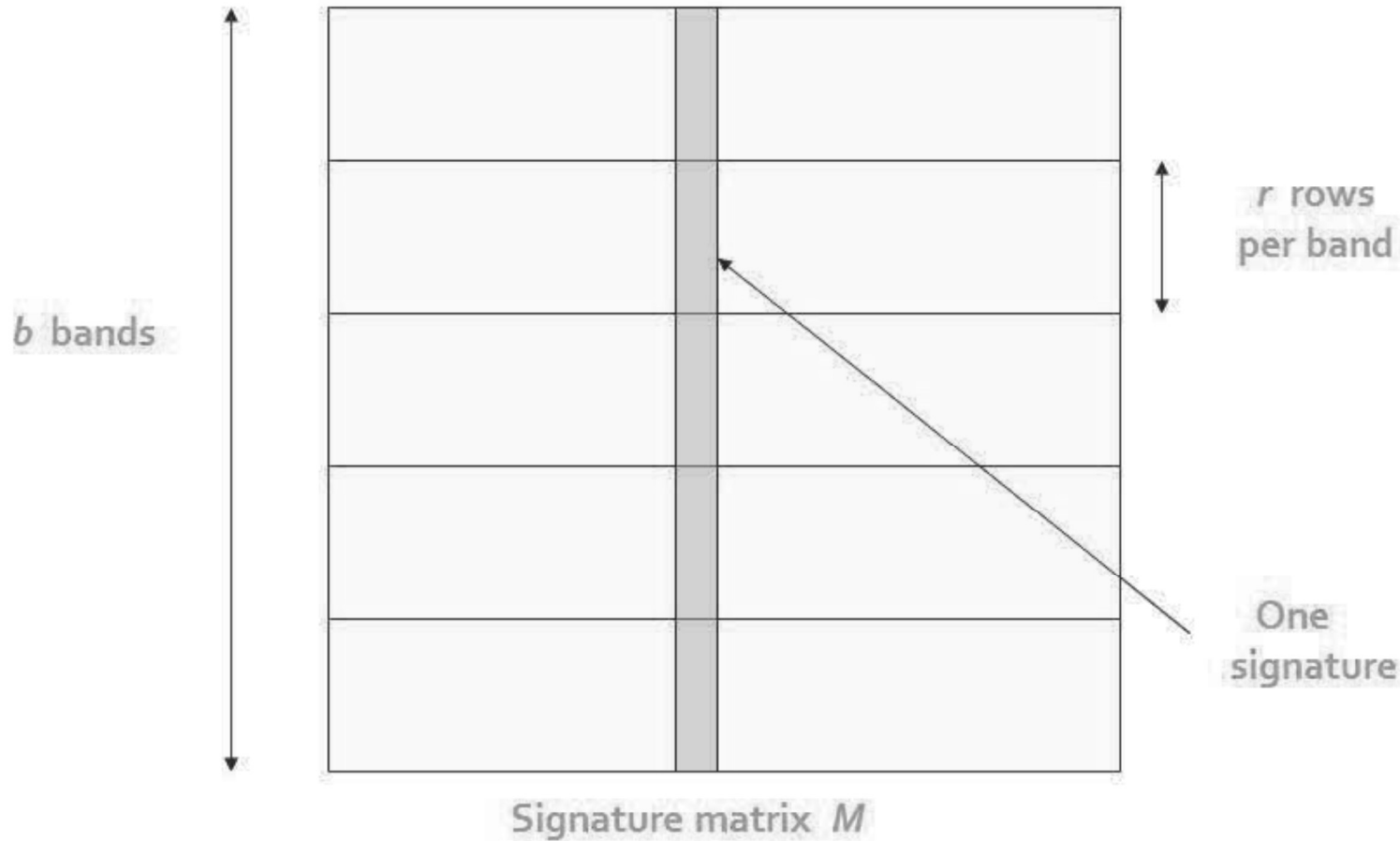
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Approach from MinHash: Hash columns of signature matrix

➡ Candidate pairs end up in the same bucket.

Locality-Sensitive Hashing

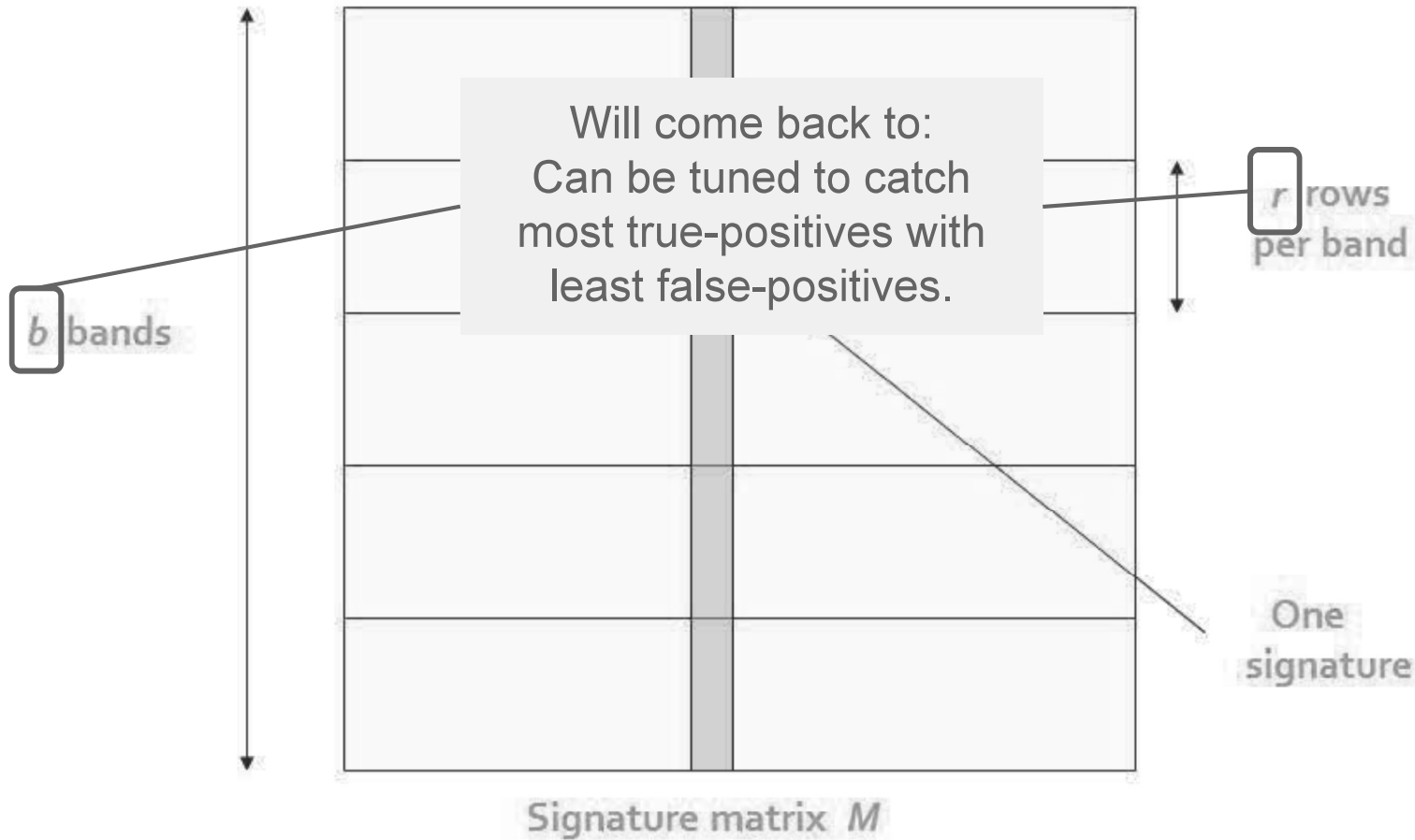
Step 1: Divide signature matrix into b bands



(Leskovec et al., 2014; <http://www.mmids.org/>)

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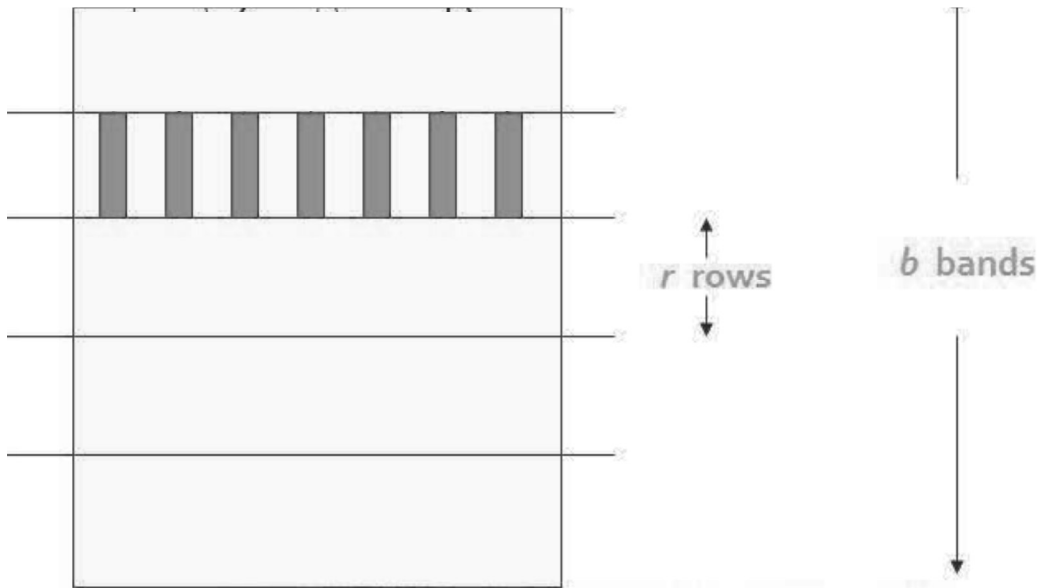
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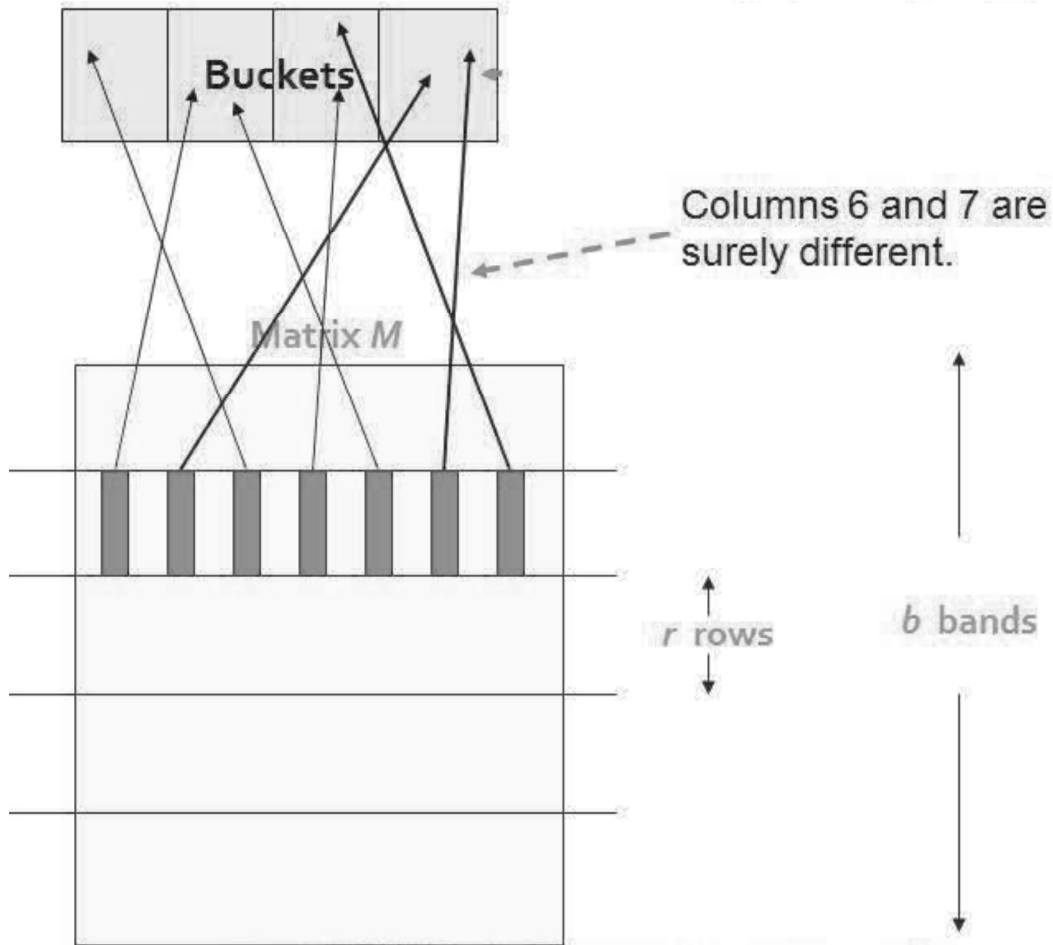
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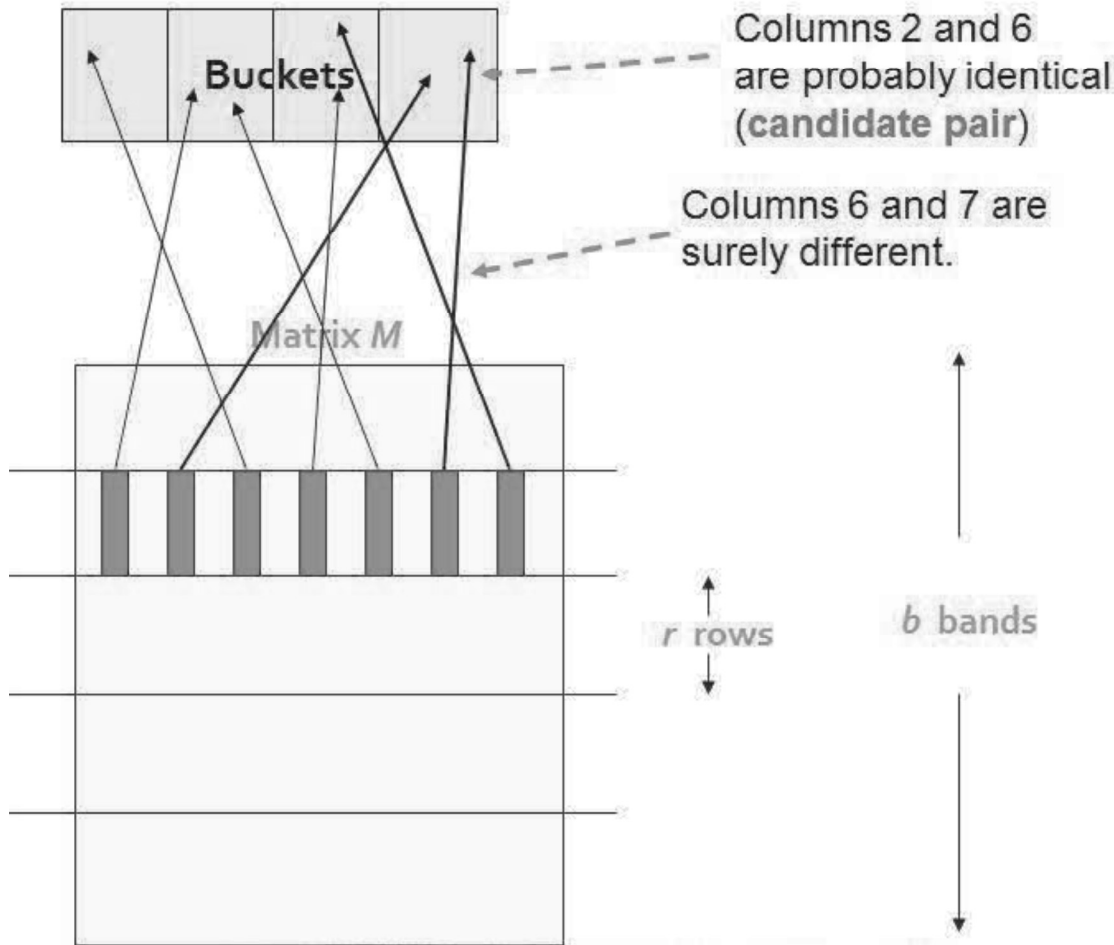
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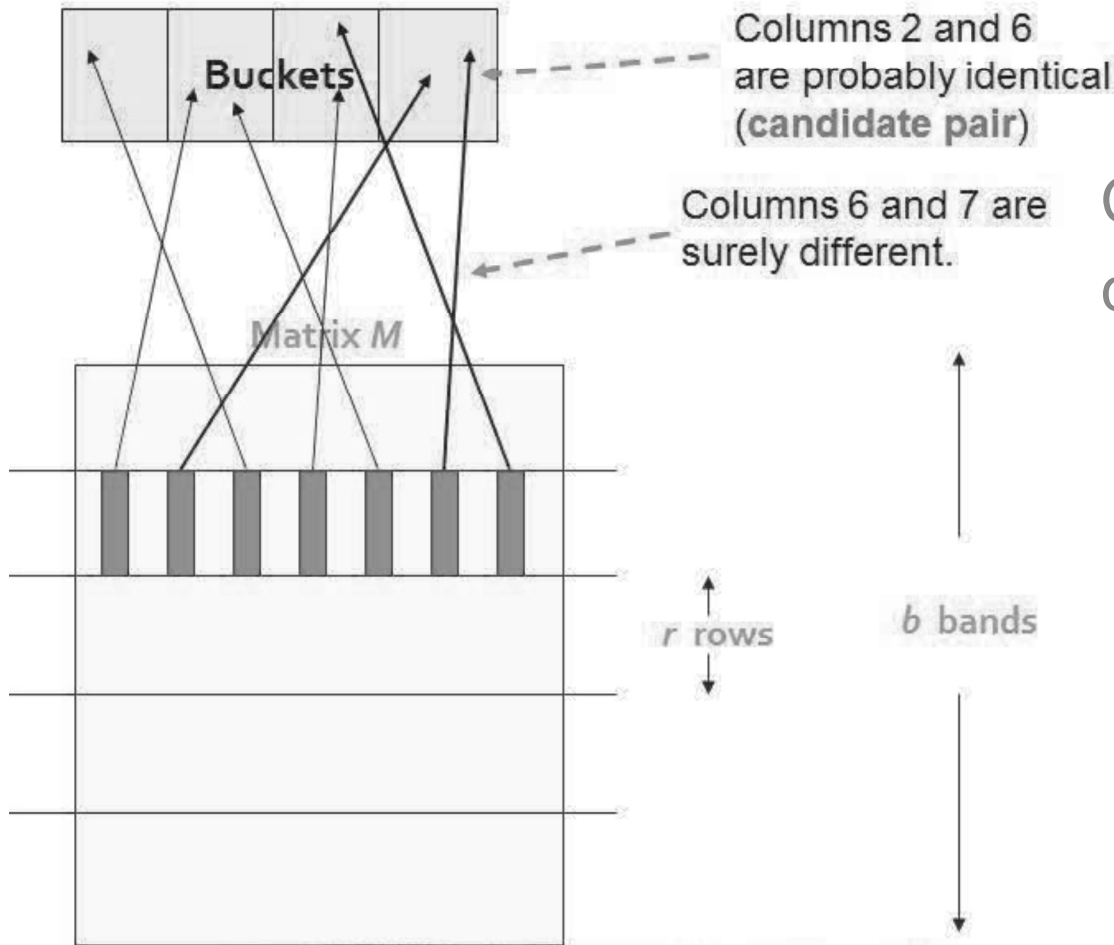
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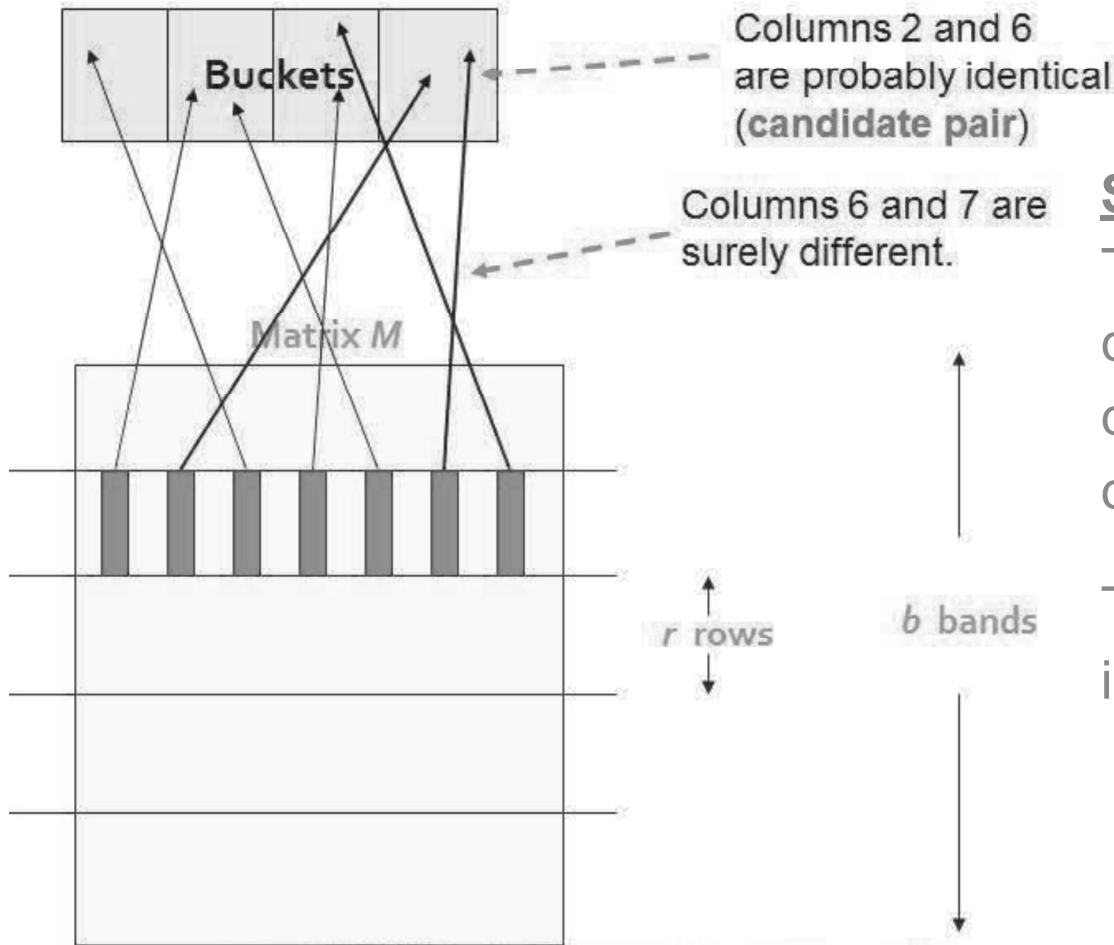


Criteria for being candidate pair:

- They end up in same bucket for at least 1 band.

Locality-Sensitive Hashing

Step 1: Divide into b bands
Step 2: Hash columns
within bands
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Simplification:

There are enough buckets compared to rows per band that columns must be identical in order to hash into same bucket.

Thus, we only need to check if identical within a band.

Document-Similarity Pipeline



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- 100,000 documents
- 100 random permutations/hash functions/rows
 - => if 4byte integers then 40Mb to hold signature matrix
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What if wanting 40% Jaccard Similarity?

Document-Similarity Pipeline

